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Carbon Capture Simulation Initiative

Tight Reformulation of Transshipment Model for Heat Integration Problems

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Introduction

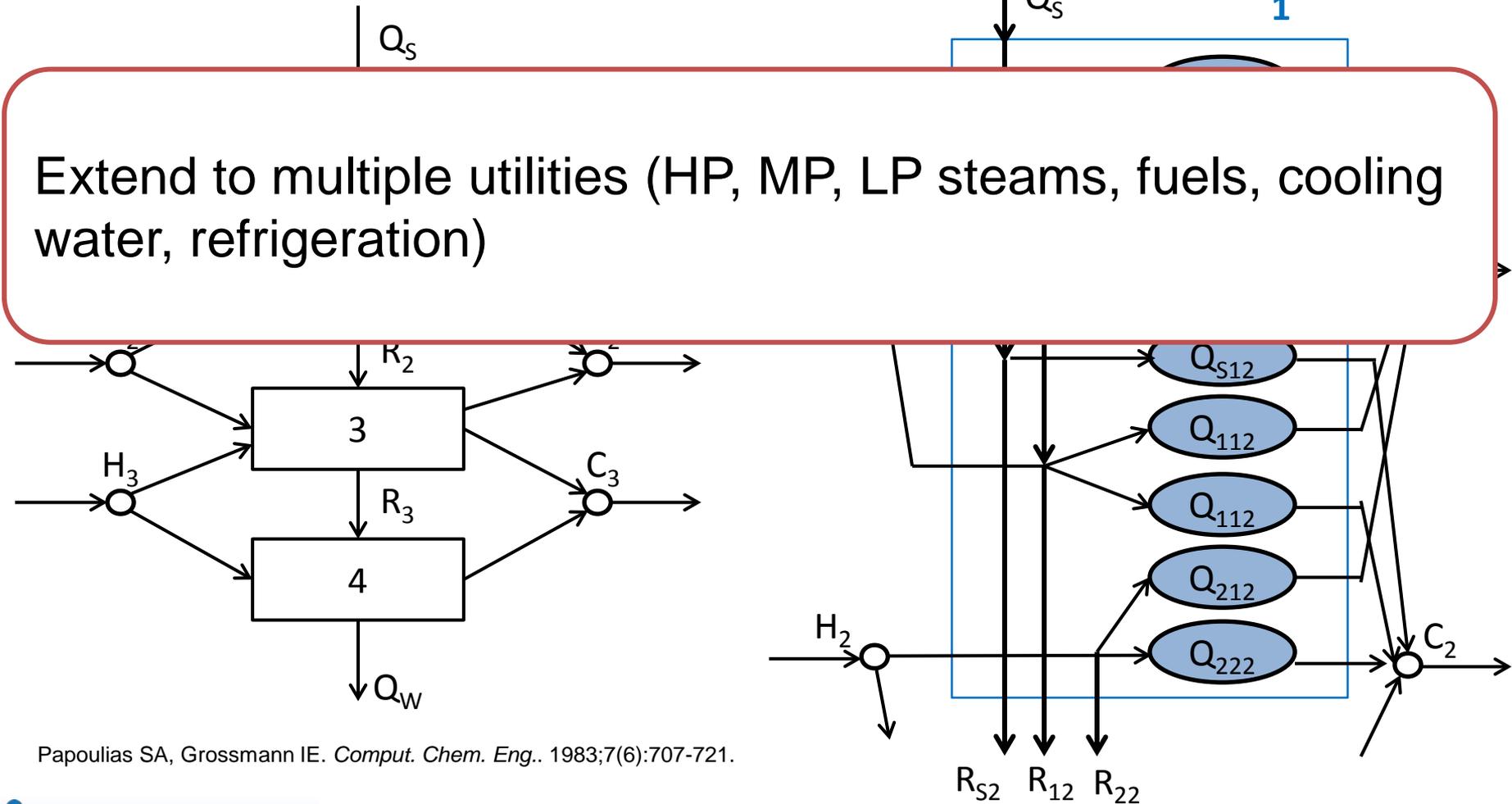
- Heat integration plays an important role to reduce energy consumption and CO₂ emissions.
- Sequential procedures to synthesize heat exchanger networks:
 - **Step 1:** Minimize utility cost → LP Transshipment Model
 - **Step 2:** Predict optimal stream matches for minimizing the number of heat exchangers → MILP Transshipment Model
 - **Step 3:** Derive heat exchanger network structures for minimizing the investment cost → NLP Model
- Sequential approach is a practical way to solve large scale heat integration problems.

Goal: Study alternative approaches for solving MILP Transshipment model.

Transshipment Model

Compact

Expanded



Papoulias SA, Grossmann IE. *Comput. Chem. Eng.*. 1983;7(6):707-721.

Transshipment Model Formulations

LP Transshipment Model

$$\min Z = \sum_{m \in S} c_m Q_m^S + \sum_{n \in W} c_n Q_n^W$$

$$\text{s.t. } R_{ik} - R_{i,k-1} + \sum_{j \in C_k} Q_{ijk} + \sum_{n \in W_k} Q_{ink} = Q_{ik}^H \quad i \in H'_k$$

$$R_{mk} - R_{m,k-1} + \sum_{j \in C_k} Q_{mjk} - Q_m^S = 0 \quad m \in S'_k$$

$$\sum_{i \in H_k} Q_{ijk} + \sum_{m \in S_k} Q_{mjk} = Q_{jk}^C \quad j \in C_k$$

$$\sum_{i \in H_k} Q_{ink} - Q_n^W = 0 \quad n \in W_k \quad k = 1, \dots, K$$

$$R_{ik}, R_{mk}, Q_{ijk}, Q_{mjk}, Q_{ink}, Q_m^S, Q_n^W \geq 0 \quad R_{i0} = R_{iK} = 0$$

Heat Balances

| | |
|-------|----------------------------------|
| Q^S | heat load of hot utility |
| Q^W | heat load of cold utility |
| Q^H | heat load of hot process stream |
| Q^C | heat load of cold process stream |
| Q | exchange of heat |
| R | heat residual |
| c | unit cost of utility |
| k | temperature interval |
| i | hot process stream |
| j | cold process stream |
| m | hot utility |
| n | cold utility |

MILP Transshipment Model

$$\min \sum_{i \in H} \sum_{j \in C} y_{ij}^p$$

$$\text{s.t. } R_{ik} - R_{i,k-1} + \sum_{j \in C_k} Q_{ijk} = Q_{ik}^H \quad i \in H'_k$$

$$\sum_{i \in H_k} Q_{ijk} = Q_{jk}^C \quad j \in C_k \quad k = 1, \dots, K_q$$

$$\sum_{k=1}^{K_p} Q_{ijk} - U_{ij} y_{ij}^p \leq 0$$

$$R_{ik}, Q_{ijk} \geq 0$$

Heat Balances

Match Constraints

| | |
|-----|--------------------------|
| y | stream match |
| Q | exchange of heat |
| R | heat residual |
| U | upper bound of heat load |
| p | subnetwork |
| k | temperature interval |
| i | hot stream |
| j | cold stream |

MILP Transshipment model is difficult to solve

- Computational time increases exponentially with the problem size.

Balanced Streams

| Case | Solution | CPU Time (s) |
|----------|----------|--------------|
| 5H, 5C | 24 | 0.5 |
| 8H, 8C | 35 | 35.9 |
| 10H, 10C | 42 | 1017.9 |
| 12H, 12C | 48 | 68688.6 |
| 15H, 15C | 57 | > 100000 |

Absolute Gap = 0.99

FCp: 0.8 ~ 2.8 MW/°C

Unbalanced Streams

| Case | Solution | CPU Time (s) |
|----------|----------|--------------|
| 5H, 5C | 26 | 0.3 |
| 10H, 10C | 39 | 25.7 |
| 15H, 15C | 55 | 660.1 |
| 17H, 17C | 67 | > 100000 |
| 20H, 20C | 78 | > 100000 |

Absolute Gap = 0.99

FCp: 0.1 ~ 14 MW/°C

Somewhat easier to solve!

- Reasons for slow computational speed:
 - Large LP relaxation gap
 - Unit coefficients in the objective function

Approaches to Reduce Computation

Model Reformulation

- Disaggregated Models
- Additional Integer Cuts
- Priority for Integer Variables

Model Modification

- Weighted Factors in Objective Function

Approximate Approaches

- Relative Optimality Gap
- Heuristic for Reduced MILP Model
- NLP Reformulation

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Disaggregated Models

Original Transshipment Model

$$\sum_{k=1}^{K_p} Q_{ijk} - U_{ij} y_{ij}^p \leq 0, \quad \forall i \in H, j \in C$$

$$U_{ij} = \min \left\{ \sum_{k=1}^{K_p} Q_{ik}^H, \sum_{k=1}^{K_p} Q_{jk}^C \right\}$$

Disaggregated Transshipment Model

$$Q_{ijk} - U_{ijk} y_{ij}^p \leq 0, \quad \forall i \in H, j \in C, k \in K$$

$$U_{ijk} = \min \left\{ \sum_{l=1}^k Q_{il}^H, Q_{jk}^C \right\}$$

Transportation Model

$$q_{iljk} - U_{iljk} y_{ij}^p \leq 0, \quad \forall i \in H, j \in C, l \in K, k \in K$$

$$U_{iljk} = \min \{ Q_{il}^H, Q_{jk}^C \}$$

$$\sum_{l=1}^k q_{iljk} = Q_{ijk}$$

- y stream match
- Q, q exchange of heat
- Q^H heat load of hot process stream
- Q^C heat load of cold process stream
- U upper bound of heat load
- p subnetwork
- k, l temperature interval
- i hot stream
- j cold stream

LP Relaxations of Disaggregated Models

| Case | Solution | LP Relaxation | | |
|----------------------------------|----------|------------------------------|-----------------------------------|----------------------|
| | | Original Transshipment Model | Disaggregated Transshipment Model | Transportation Model |
| <i>Balanced Streams</i> | | | | |
| 5H, 5C | 24 | 16.302 | 16.718 | 16.802 |
| 8H, 8C | 35 | 24.357 | 24.755 | 24.791 |
| 10H, 10C | 42 | 28.848 | 30.075 | 30.103 |
| 12H, 12C | 48 | 32.135 | 33.395 | 33.629 |
| 15H, 15C | 57 | 40.390 | 42.388 | 42.576 |
| <i>Unbalanced Streams</i> | | | | |
| 5H, 5C | 26 | 17.931 | 20.072 | 20.645 |
| 10H, 10C | 39 | 29.969 | 31.899 | 32.609 |
| 15H, 15C | 55 | 41.424 | 43.477 | 44.640 |
| 17H, 17C | 67 | 48.839 | 52.636 | 53.551 |
| 20H, 20C | 78 | 56.593 | 61.848 | 63.231 |

Disaggregated models: tighter LP relaxations.

Computational Performance of Disaggregated Models

| Case | Solution | CPU Time (s) | | |
|---------------------------|----------|------------------------------|-----------------------------------|----------------------|
| | | Original Transshipment Model | Disaggregated Transshipment Model | Transportation Model |
| Balanced Streams | | | | |
| 5H, 5C | 24 | 0.5 | 0.5 | 0.4 |
| 8H, 8C | 35 | 35.9 | 34.9 | 91.1 |
| 10H, 10C | 42 | 1017.9 | 1011.4 | 3075.1 |
| 12H, 12C | 48 | 68688.6 | 36356.6 | > 100000 |
| 15H, 15C | 57 | > 100000 | > 100000 | > 100000 |
| Unbalanced Streams | | | | |
| 5H, 5C | 26 | 0.3 | 0.2 | 0.4 |
| 10H, 10C | 39 | 25.7 | 21.1 | 150.1 |
| 15H, 15C | 55 | 660.1 | 1043.1 | > 100000 |
| 17H, 17C | 67 | > 100000 | 76676.2 | > 100000 |
| 20H, 20C | 78 | > 100000 | > 100000 | > 100000 |

Transshipment MILP model outperforms Transportation MILP model.

Disaggregated Transshipment model: shorter computational times than the original model in most of cases.

Additional Integer Cuts

$$\sum_{j \in C} y_{ij}^p \geq \left\lceil \frac{\sum_{k=1}^{K_p} Q_{ik}^H}{\max_{j \in C} \left\{ \sum_{k=1}^{K_p} Q_{jk}^C \right\}} \right\rceil, \forall i \in H \quad \sum_{i \in H} y_{ij}^p \geq \left\lceil \frac{\sum_{k=1}^{K_p} Q_{jk}^C}{\max_{i \in H} \left\{ \sum_{k=1}^{K_p} Q_{ik}^H \right\}} \right\rceil, \forall j \in C \quad \sum_{i \in H} \sum_{j \in C} y_{ij}^p \leq \sum_{i \in H} N_i^p + \sum_{j \in C} N_j^p$$

y stream match
 Q^H, Q^C heat load of hot/cold stream
 p subnetwork
 k temperature interval
 i hot stream
 j cold stream
 N number of streams

LP Relaxation

| Case | Solution | LP Relaxation | | | |
|---------------------------|----------|------------------------------|---------|-----------------------------------|---------|
| | | Original Transshipment Model | | Disaggregated Transshipment Model | |
| | | w/o cuts | w/ cuts | w/o cuts | w/ cuts |
| Balanced Streams | | | | | |
| 5H, 5C | 24 | 16.302 | 17.383 | 16.718 | 17.642 |
| 8H, 8C | 35 | 24.357 | 24.416 | 24.755 | 24.850 |
| 10H, 10C | 42 | 28.848 | 29.852 | 30.075 | 30.876 |
| 12H, 12C | 48 | 32.135 | 32.854 | 33.395 | 34.268 |
| 15H, 15C | 57 | 40.390 | 40.633 | 42.388 | 42.662 |
| Unbalanced Streams | | | | | |
| 5H, 5C | 26 | 17.931 | 18.178 | 20.072 | 20.645 |
| 10H, 10C | 39 | 29.969 | 30.067 | 31.899 | 32.609 |
| 15H, 15C | 55 | 41.424 | 41.604 | 43.477 | 44.669 |
| 17H, 17C | 67 | 48.839 | 49.246 | 52.636 | 53.674 |
| 20H, 20C | 78 | 56.593 | 56.816 | 61.848 | 63.301 |

Computational Performance with Additional Integer Cuts

| Case | Solution | CPU Time (s) | | | |
|---------------------------|----------|------------------------------|----------|-----------------------------------|-----------|
| | | Original Transshipment Model | | Disaggregated Transshipment Model | |
| | | w/o cuts | w/ cuts | w/o cuts | w/ cuts |
| Balanced Streams | | | | | |
| 5H, 5C | 24 | 0.5 | 0.5 | 0.5 | 0.5 |
| 8H, 8C | 35 | 35.9 | 33.0 | 34.9 | 33.2 |
| 10H, 10C | 42 | 1017.9 | 1039.9 | 1011.4 | 878.3 |
| 12H, 12C | 48 | 68688.6 | 65506.2 | 36356.6 | 33869.2 |
| 15H, 15C | 57 | > 100000 | > 100000 | > 100000 | > 100000 |
| Unbalanced Streams | | | | | |
| 5H, 5C | 26 | 0.3 | 0.3 | 0.2 | 0.3 |
| 10H, 10C | 39 | 25.7 | 16.5 | 21.1 | 30.7 |
| 15H, 15C | 55 | 660.1 | 919.2 | 1043.1 | 749.8 |
| 17H, 17C | 67 | > 100000 | > 100000 | 76676.2 | 28682.359 |
| 20H, 20C | 78 | > 100000 | > 100000 | > 100000 | > 100000 |

Integer cuts: increase computational speed in most of cases.

Disaggregated Transshipment model: the best MILP formulation.

Branching Priority for Binary Variables

$$y_{ij}.\text{prior} = 1/\text{UB}_{ij}$$

Branch y_{ij} with largest upper bound (UB_{ij}) first

| Case | Solution | CPU Time (s) | | | |
|---------------------------|----------|------------------------------|-------------|-----------------------------------|-------------|
| | | Original Transshipment Model | | Disaggregated Transshipment Model | |
| | | w/o priority | w/ priority | w/o priority | w/ priority |
| Balanced Streams | | | | | |
| 5H, 5C | 24 | 0.5 | 0.5 | 0.5 | 0.5 |
| 8H, 8C | 35 | 33.0 | 27.6 | 33.2 | 29.6 |
| 10H, 10C | 42 | 1039.9 | 680.0 | 878.3 | 607.0 |
| 12H, 12C | 48 | 65506.2 | 39856.9 | 33869.2 | 24400.8 |
| 15H, 15C | 57 | > 100000 | > 100000 | > 100000 | > 100000 |
| Unbalanced Streams | | | | | |
| 5H, 5C | 26 | 0.3 | 0.4 | 0.3 | 0.3 |
| 10H, 10C | 39 | 16.5 | 6.9 | 30.7 | 31.3 |
| 15H, 15C | 55 | 919.2 | 648.2 | 749.8 | 1527.2 |
| 17H, 17C | 67 | > 100000 | > 100000 | 28682.359 | > 100000 |
| 20H, 20C | 78 | > 100000 | > 100000 | > 100000 | > 100000 |

Branching priority: improves the performance in most of cases.

Disaggregated Transshipment model with additional integer cuts and branching priority : the base model in the following studies.



Approaches to Reduce Computation

Model Reformulation

- Disaggregated Models
- Additional Integer Cuts
- Priority for Integer Variables

Model Modification

- Weighted Factors in Objective Function

Approximate Approaches

- Relative Optimality Gap
- Heuristic for Reduced MILP Model
- NLP Reformulation

Weighted Factors in Objective Function

Objective Function:

$$\min \sum_{i \in H} \sum_{j \in C} y_{ij}$$



$$\min \sum_{i \in H} \sum_{j \in C} w_{ij} y_{ij}$$

- i hot stream
- j cold stream
- y stream match
- w weighting factor
- U upper bound of heat load
- ΔT mean temperature difference
- ΔT_{in} inlet temperature difference
- ΔT_{out} outlet temperature difference

$$w_{ij} = \frac{U_{ij}}{\Delta T_{ij}}$$

$$\Delta T_{ij} = \frac{\Delta T_{out,ij} - \Delta T_{in,ij}}{\ln \frac{\Delta T_{out,ij}}{\Delta T_{in,ij}}}$$

$$\Delta T_{ij} \cong \left(\Delta T_{in,ij} \Delta T_{out,ij} \frac{\Delta T_{in,ij} + \Delta T_{out,ij}}{2} \right)^{1/3}$$

Computational Performance of Weighted Model

| Case | Base Model | | Weighted Model | |
|---------------------------|------------|--------------|----------------|--------------|
| | Solution | CPU Time (s) | Solution | CPU Time (s) |
| Balanced Streams | | | | |
| 5H, 5C | 24 | 0.5 | 25 | 0.4 |
| 8H, 8C | 35 | 29.6 | 38 | 22.9 |
| 10H, 10C | 42 | 607.0 | 47 | 642.8 |
| 12H, 12C | 48 | 24400.8 | 53 | 18608.8 |
| 15H, 15C | 57 | > 100000 | 65 | > 100000 |
| Unbalanced Streams | | | | |
| 5H, 5C | 26 | 0.3 | 26 | 0.2 |
| 10H, 10C | 39 | 31.3 | 43 | 3.2 |
| 15H, 15C | 55 | 1527.2 | 61 | 145.4 |
| 17H, 17C | 67 | > 100000 | 75 | 1901.0 |
| 20H, 20C | 78 | > 100000 | 85 | > 100000 |

Absolute Gap = 0.99

Relative Gap = 1%

Investment Cost of Weighted Model

Minimize the **investment cost** of heat exchanger networks for both models in SYNHEAT, by fixing all stream matches in previous results.

| Case | Base Model | | | Weighted Model | | |
|----------------------------------|----------------------|------------------------------|-------------------------|----------------------|------------------------------|-------------------------|
| | # of Heat Exchangers | Total Area (m ²) | Investment Cost (\$/yr) | # of Heat Exchangers | Total Area (m ²) | Investment Cost (\$/yr) |
| <i>Balanced Streams</i> | | | | | | |
| 5H, 5C | 24 | 250.6 | 67330 | 25 | 179.0 | 61630 |
| <i>Unbalanced Streams</i> | | | | | | |
| 5H, 5C | 26 | 672.9 | 141613 | 26 | 595.3 | 124451 |

Weighted model: smaller exchanger area and investment cost, but possibly more stream matches.

Approaches to Reduce Computation

Model Reformulation

- Disaggregated Models
- Additional Integer Cuts
- Priority for Integer Variables

Model Modification

- Weighted Factors in Objective Function

Approximate Approaches

- Relative Optimality Gap
- Heuristic for Reduced MILP Model
- NLP Reformulation

Different Optimality Gap

| Case | Absolute Gap = 0.99 (Base) | | Relative Gap = 5% | | Relative Gap = 10% | |
|---------------------------|----------------------------|--------------|-------------------|--------------|--------------------|--------------|
| | Solution | CPU Time (s) | Solution | CPU Time (s) | Solution | CPU Time (s) |
| Balanced Streams | | | | | | |
| 5H, 5C | 24 | 0.5 | 24 | 0.5 | 24 | 0.5 |
| 8H, 8C | 35 | 29.6 | 35 | 29.6 | 35 | 18.3 |
| 10H, 10C | 42 | 607.0 | 42 | 607.0 | 42 | 187.5 |
| 12H, 12C | 48 | 24400.8 | 48 | 24400.8 | 48 | 3877.9 |
| 15H, 15C | 57 | > 100000 | 57 | > 100000 | 57 | > 100000 |
| Unbalanced Streams | | | | | | |
| 5H, 5C | 26 | 0.3 | 26 | 0.3 | 26 | 0.3 |
| 10H, 10C | 39 | 31.3 | 39 | 31.3 | 40 | 4.5 |
| 15H, 15C | 55 | 1527.2 | 55 | 1442.9 | 56 | 109.8 |
| 17H, 17C | 67 | > 100000 | 67 | >100000 | 69 | 6414.2 |
| 20H, 20C | 78 | > 100000 | 78 | >100000 | 78 | 30682.8 |

10% of relative gap is acceptable for most of cases.

Heuristic Approach to Solve Reduce MILP Model

- **Step 1:** Solve the relaxed LP model (Transshipment or Transportation).
- **Step 2:** Fix $y_{ij} = 0$ in the full MILP model if its relaxed value = 0 in the LP model. The set for these y_{ij} is defined as Y_0^{rx} . Solve the reduced MILP model (Disaggregated Transshipment).
- **Step 3:** In the full MILP model, fix all $y_{ij} \notin Y_0^{rx}$ as the solution of the reduced MILP model.
- **Step 4:** Fix one of $y_{ij} \in Y_0^{rx}$ to be 1, and leave other $y_{ij} \in Y_0^{rx}$ as variables. Solve the test MILP problem.
- **Step 5:** Check the value of Q_{ij} for the above y_{ij} . If $Q_{ij} = 0$, keep y_{ij} in Y_0^{rx} . If $Q_{ij} > 0$, remove y_{ij} from Y_0^{rx} .
- **Step 6:** Relax the above y_{ij} to variable. Go to the next y_{ij} .
- **Repeat Step 4 - 6** for every $y_{ij} \in Y_0^{rx}$.
- **Step 7:** Fix $y_{ij} = 0$ for all $y_{ij} \in Y_0^{rx}$ in the full MILP model. Solve the final reduced MILP model (Disaggregated Transshipment) and obtain the approximate solution.

Heuristic for Reduced MILP Model

| Case | Full Model (Base) | | Reduced Model by LP Relaxation of Original Transshipment Model | | Reduced Model by LP Relaxation of Disaggregated Transshipment Model | | Reduced Model by LP Relaxation of Transportation Model | |
|---------------------------|-------------------|--------------|--|--------------|--|--------------|--|--------------|
| | Solution | CPU Time (s) | Solution | CPU Time (s) | Solution | CPU Time (s) | Solution | CPU Time (s) |
| Balanced Streams | | | | | | | | |
| 5H, 5C | 24 | 0.5 | 25 | 2.8 | 24 | 2.9 | 24 | 2.8 |
| 8H, 8C | 35 | 29.6 | 35 | 8.6 | 35 | 6.4 | 36 | 6.6 |
| 10H, 10C | 42 | 607.0 | 44 | 8.0 | 43 | 92.0 | 43 | 68.8 |
| 12H, 12C | 48 | 24400.8 | 51 | 168.7 | 48 | 348.2 | 48 | 1121.7 |
| 15H, 15C | 57 | > 100000 | 58 | > 100000 | 59 | > 100000 | 59 | > 100000 |
| Unbalanced Streams | | | | | | | | |
| 5H, 5C | 26 | 0.3 | 26 | 2.4 | 26 | 2.4 | 26 | 2.7 |
| 10H, 10C | 39 | 31.3 | 42 | 4.3 | 41 | 4.5 | 41 | 5.9 |
| 15H, 15C | 55 | 1527.2 | 61 | 9.9 | 56 | 147.0 | 55 | 521.9 |
| 17H, 17C | 67 | > 100000 | 72 | 52.2 | 70 | 23645.9 | 70 | 10157.0 |
| 20H, 20C | 78 | > 100000 | 86 | 63.9 | 79 | > 100000 | 82 | > 100000 |

**Reduced MILP models:
Good approximate solutions
CPU times reduced by at least one order of magnitude**

NLP Reformulation

$$y_{ij} \in \{0, 1\} \longrightarrow y_{ij} \in [0, 1]$$

~~Option 1:~~

~~Adding to Constraints:~~

~~$$y_{ij}(1 - y_{ij}) \leq \epsilon, \quad \forall i \in H, j \in C$$~~

Option 2:

Objective Function:

$$\sum_{i \in H} \sum_{j \in C} y_{ij} + K \sum_{i \in H} \sum_{j \in C} y_{ij}(1 - y_{ij})$$

Computational Performance of NLP Model

| Case | MILP Model (Base) | | NLP Model | |
|---------------------------|-------------------|--------------|-----------|--------------|
| | Solution | CPU Time (s) | Solution | CPU Time (s) |
| Balanced Streams | | | | |
| 5H, 5C | 24 | 0.5 | 26 | 82.0 |
| 8H, 8C | 35 | 29.6 | 39 | 233.6 |
| 10H, 10C | 42 | 607.0 | 48 | 202.5 |
| 12H, 12C | 48 | 24400.8 | 53 | 458.5 |
| 15H, 15C | 57 | > 100000 | 70 | 1291.9 |
| Unbalanced Streams | | | | |
| 5H, 5C | 26 | 0.3 | 26 | 81.9 |
| 10H, 10C | 39 | 31.3 | 46 | 864.8 |
| 15H, 15C | 55 | 1527.2 | 66 | 1242.7 |
| 17H, 17C | 67 | > 100000 | 79 | 1674.1 |
| 20H, 20C | 78 | > 100000 | 95 | 14804.0 |

NLP Solver: OQNLP

NLP Models: Faster but worse solutions

Conclusions

- MILP Transshipment model is computational expensive for large-scale problems.
- Disaggregate Transshipment model is the best formulation for most of case studies.
- Additional integer cuts and branching priority for binary variables are helpful to improve computational performance.
- Weighted factors can be added into the objective function to reduce solution times and obtain designs with lower investment costs.
- An appropriate relative optimality gap (10%) can be used to get near-optimal solutions in relatively short times.
- Reduced MILP models obtained by a heuristic approach achieves near-optimal solutions while reducing solution times by at least one order of magnitude.
- NLP formulations obtained approximate solutions in relatively short times.

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DOE: Carbon Capture Simulation Initiative (CCSI)

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